

# Signomial and Polynomial Optimization via Relative Entropy and Partial Dualization

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Joint work with Venkat Chandrasekaran and Adam Wierman.

# Definitions from convex analysis



$X \subset \mathbb{R}^n$  is a **convex set** if it contains all of its line segments.

$f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a **convex function** if

$$f(t\mathbf{x} + (1 - t)\mathbf{y}) \leq tf(\mathbf{x}) + (1 - t)f(\mathbf{y}).$$

for all  $\mathbf{x}, \mathbf{y} \in \text{dom } f$  and all  $t \in [0, 1]$ .

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The **relative entropy function** continuously extends

$$D(\mathbf{u}, \mathbf{v}) = \sum_{i=1}^m u_i \log(u_i/v_i) \quad \text{to} \quad \mathbb{R}_+^m \times \mathbb{R}_+^m.$$

# Convex duality



Start with a **primal** problem

$$\text{Val}(\mathbf{c}) = \inf_{\mathbf{x}} \{\mathbf{c}^\top \mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}.$$

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$$\text{Val}(\mathbf{c}) \geq L.$$

Write such a constraint as: *there exists a  $\boldsymbol{\mu}$  where*

$$\mathbf{A}^\top \boldsymbol{\mu} \leq \mathbf{c} \quad \text{and} \quad \mathbf{b}^\top \boldsymbol{\mu} \geq L.$$



# Nonnegativity and Optimization



Start with a function  $f$  and a set  $X \subset \mathbb{R}^n$

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- 3 Find largest  $\gamma$  so  $c(\gamma)$  belongs to  $\tilde{C}(\phi, X)$ .

# Our functions of interest

**Caltech***polynomials**signomials*

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Countable basis.

Complexity measured by degree

$$\max_i \alpha_{i1} + \cdots + \alpha_{in}.$$

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Uncountable basis.

Complexity measured by

number of terms  $m$ .

# How we think of signomials



Signomials are often written  $\mathbf{y} \mapsto \sum_{i=1}^m c_i \mathbf{y}^{\alpha_i}$ , with  $\mathbf{y} \in \mathbb{R}_{++}^n$ .

The *exponential form* has a powerful connection to convexity.

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One use of your existing intuition:

- Pick an “interesting” polynomial  $p$ .
- Define  $f(\mathbf{z}) = p(\exp(\mathbf{z}) - \exp(-\mathbf{z}))$ .
- $f$  will behave similarly to  $p$ , near  $\mathbf{0}$ .

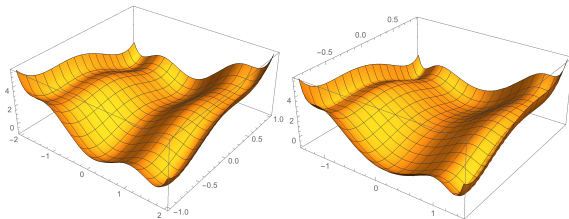
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- Simple “SAGE relaxations.”
- Partial dualization.
- Two examples.



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Definition → Representation → Example.

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## 3 SAGE polynomials

Definition → Example → Representation.

# The signomial $X$ -nonnegativity cones



The  $X$ -nonnegativity cone for signomials over exponents  $\alpha$ :

$$C_{\text{NNS}}(\alpha, X) \doteq \left\{ \mathbf{c} : \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}) \geq 0 \text{ for all } \mathbf{x} \text{ in } X \right\}.$$

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If

$$f(\mathbf{x}) = c_1 \exp(\mathbf{0} \cdot \mathbf{x}) + c_2 \exp(\alpha_2 \cdot \mathbf{x}) + \cdots + c_m \exp(\alpha_m \cdot \mathbf{x}),$$

then

$$\inf_{\mathbf{x} \in X} f(\mathbf{x}) = \sup \{ \gamma : \mathbf{c} - (\gamma, 0, \dots, 0) \in C_{\text{NNS}}(\alpha, X) \}.$$

# $X$ -SAGE $\Rightarrow$ $X$ -nonnegativity

*Definition.* A signomial which is nonnegative over  $X$  and which has at most one negative coefficient is an “ $X$ -AGE function.”

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*Crucial question:* How to represent the AGE cones?

# The convex duality behind AGE cones



Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

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Divide out the problematic exponential, and rearrange terms:

$$\begin{aligned} \sum_{i=1}^m c_i \exp(\alpha_i \cdot x) \geq 0 &\Leftrightarrow \sum_{i=1}^m c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq 0 \\ &\sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq -c_k. \end{aligned}$$



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The nonnegativity condition

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$$\begin{aligned} [1\alpha_k - \alpha_{\setminus k}]^\top \nu &= \lambda \quad \text{and} \\ \sigma_X(\lambda) + D(\nu, c_{\setminus k}) - \nu^\top \mathbf{1} &\leq c_k. \end{aligned}$$

# Tractability of $X$ -SAGE cones

There are two constraints in an AGE cone:

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When does it have a closed form? Examples include ...

$$X = \mathbb{R}^n \quad \Rightarrow \quad \sigma_X(\boldsymbol{\lambda}) = \begin{cases} 0 & \text{if } \boldsymbol{\lambda} = \mathbf{0} \\ +\infty & \text{if } \boldsymbol{\lambda} \neq \mathbf{0} \end{cases}$$

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$$X = \{\mathbf{x} : \|\mathbf{x} - \mathbf{a}\| \leq r\} \quad \Rightarrow \quad \sigma_X(\boldsymbol{\lambda}) = \boldsymbol{\lambda}^\top \mathbf{a} + r\|\boldsymbol{\lambda}\|_*$$

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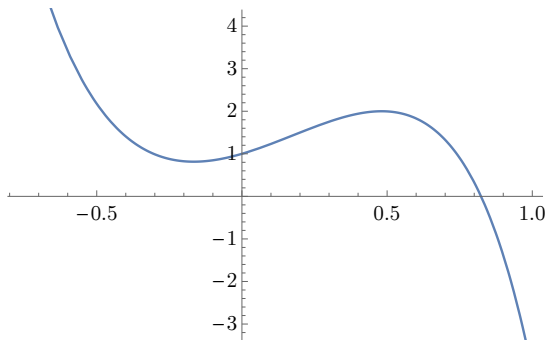
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When is it *tractable*?

**Whenever  $X$  is tractable.**

# A univariate example

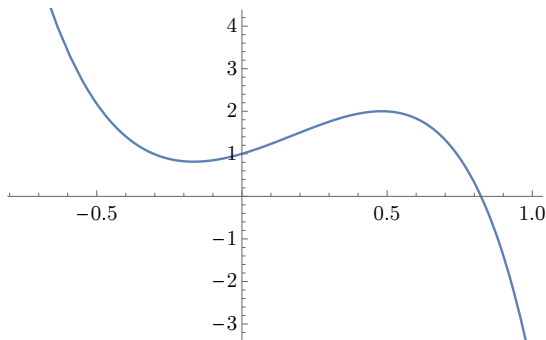
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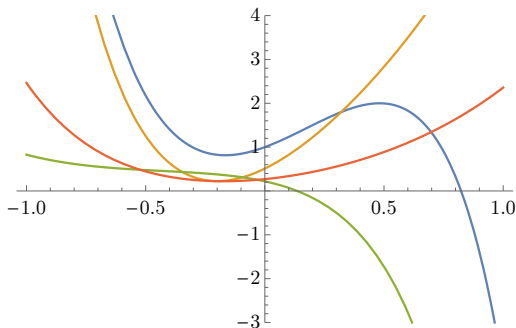
$$f_1(x) = 0.88 \cdot e^{-3x} + 0.82 \cdot e^{-2x} + 2.69 \cdot e^x + 0.12 \cdot e^{2x} - 4 \cdot e^{-x}$$

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# Optimization.

# Simple SAGE relaxations

Consider  $f(\mathbf{x}) = \sum_{i=1}^m c_i \exp(\boldsymbol{\alpha}_i \cdot \mathbf{x})$  with  $\boldsymbol{\alpha}_1 = \mathbf{0}$ .

$$\begin{aligned} \inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} &= \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{NNS}}(\boldsymbol{\alpha}, X)\} \\ &\geq \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)\} \\ &= \inf\left\{\mathbf{c}^\top \mathbf{v} : \begin{array}{l} \mathbf{v} \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)^\dagger \\ \text{satisfies } \mathbf{v} \cdot \mathbf{e}_1 = 1 \end{array}\right\} \end{aligned}$$

What about [solution recovery](#)?

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What about [solution recovery](#)? When  $X$  is convex, we have

$$\begin{aligned} C_{\text{SAGE}}(\boldsymbol{\alpha}, X)^\dagger &= \text{cl}\{\mathbf{v} : \text{some } \mathbf{z}_1, \dots, \mathbf{z}_m \text{ in } \mathbb{R}^n \text{ satisfy} \\ &\quad v_k \log(\mathbf{v}/v_k) \geq [\boldsymbol{\alpha} - \mathbf{1}\boldsymbol{\alpha}_k] \mathbf{z}_k \\ &\quad \text{and } \mathbf{z}_k/v_k \in X \text{ for all } k \text{ in } [m]\}. \end{aligned}$$

# An example in $\mathbb{R}^3$

Minimize

$$f(\mathbf{x}) = 0.5 \exp(x_1 - x_2) - \exp x_1 - 5 \exp(-x_2)$$

over

$$X = \{\mathbf{x} : \log 70 \leq x_1 \leq \log 150,$$

$$\log 1.0 \leq x_2 \leq \log 30,$$

$$\log 0.5 \leq x_3 \leq \log 21$$

$$\exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100\}.$$

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$$\text{Compute } f_X^{\text{SAGE}} = -147.85713 \leq f_X^*,$$

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over

$$\begin{aligned} X = \{ \mathbf{x} : & \log 70 \leq x_1 \leq \log 150, \\ & \log 1.0 \leq x_2 \leq \log 30, \\ & \log 0.5 \leq x_3 \leq \log 21 \\ & \exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100 \}. \end{aligned}$$

Compute  $f_X^{\text{SAGE}} = -147.85713 \leq f_X^*$ , recover feasible

$$\mathbf{x}^* = (5.01063529, 3.40119660, -0.48450710)$$

satisfying  $f(\mathbf{x}^*) = -147.66666$ .



# An example in $\mathbb{R}^3$

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satisfying  $f(\mathbf{x}^*) = -147.66666$ . *This is actually optimal!*

# Nonconvex constraints



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Then, if  $\lambda_1, \dots, \lambda_k$  are nonnegative dual variables, we have

$$\inf_{\mathbf{x} \in \Omega} f(\mathbf{x}) \geq \sup \left\{ \gamma : f + \sum_{i=1}^k \lambda_i g_i - \gamma \text{ is } X\text{-SAGE} \right\}.$$

# The SimPleAC aircraft design problem



From Warren Hoburg's PhD thesis.

Problem statistics:

- 140 variables.
- 89 inequality constraints (1 nonconvex).
- 67 equality constraints (15 nonconvex).

Performance of the most basic SAGE relaxation:

- bound "cost  $\geq 2957$ " (roughly match a known solution).
- MOSEK solves in two seconds, on a six year old laptop.
- solution recovery fails (numerical issues).

# SAGE polynomials

# $X$ -nonnegative and $X$ -AGE polynomials



Recall our standard notation  $\mathbf{x}^{\alpha_i} \doteq \prod_{j=1}^n x_j^{\alpha_{ij}}$ .

The matrix  $\alpha$  and a set  $X \subset \mathbb{R}^n$  induce a nonnegativity cone

$$C_{\text{NNP}}(\alpha, X) = \{ \mathbf{c} : c_1 \mathbf{x}^{\alpha_1} + \cdots + c_m \mathbf{x}^{\alpha_m} \geq 0 \text{ for all } \mathbf{x} \text{ in } X \}.$$



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*Def.*  $f(\mathbf{x}) = c_1 \mathbf{x}^{\alpha_1} + \cdots + c_m \mathbf{x}^{\alpha_m}$  is an  **$X$ -AGE polynomial** if

- 1  $\mathbf{c}$  belongs to  $C_{\text{NNP}}(\alpha, X)$ , and
- 2 at most one “ $k$ ” has  $c_k \mathbf{x}^{\alpha_k} < 0$  for some  $\mathbf{x} \in X$ .

# A symbolic example

Consider  $f(x) = n - \sum_{i=1}^n \prod_{j \in [n] \setminus \{i\}} x_j$  over  $X = \{-1, 1\}^n$ .

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- 2 Therefore  $f_i(\mathbf{x}) = 1 - \prod_{j \neq i} x_j$  are  $X$ -AGE.
- 3 Since  $f = \sum_{i=1}^n f_i$ , conclude  $f$  is  $X$ -SAGE.

# The representation problem

The cone of coefficients for  $X$ -**SAGE polynomials** is given by

$$C_{\text{SAGE}}^{\text{POLY}}(\alpha, X) = \sum_{k=1}^m \left\{ \mathbf{c} \mid \begin{array}{l} \mathbf{c} \text{ in } C_{\text{NNP}}(\alpha, X), \text{ and for} \\ i \neq k, \mathbf{x} \in X \Rightarrow c_i \mathbf{x}^{\alpha_i} \geq 0 \end{array} \right\}.$$

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- SAGE signomials put monomials front and center.
- Monomials “ $\mathbf{x}^{\alpha_i}$ ” are quite different from “ $\exp(\alpha_i \cdot \mathbf{x})$ ”.

## Representation: the single-orthant case

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# Representation: sign-symmetry

Consider  $X \subset \mathbb{R}^n$  which satisfies

- 1 invariance under reflection about  $\{\mathbf{x} : x_i = 0\}$ ,
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It can subsequently be shown that

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If  $\alpha_i$  isn't even, then some  $\mathbf{x}_1, \mathbf{x}_2 \in X$  satisfy  $x_1^{\alpha_i} < 0 < x_2^{\alpha_i}$ .



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Studied by Montel (1928) and Niculescu (2000), among others.

If  $g_1, \dots, g_k$  are log-log convex on a box  $B \subset \mathbb{R}_{++}^n$ , then

$\log\{(\mathbf{x}, \mathbf{t}) : \mathbf{x} \in B, \mathbf{t} \in \mathbb{R}_{++}^k, g(\mathbf{x}) \leq \mathbf{t}\} \subset \mathbb{R}^{n+k}$  is convex.

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Some tractable constraints for  $X$ -SAGE polynomials:

$$\|\mathbf{x}\|_p \leq a \quad x_j^2 = a \quad a \leq \mathbb{P}\{\mathcal{N}(0, \sigma) \geq |x|\}$$

where  $a > 0$ .

Thank you!

# Handling conditional nonnegativity



Typically, one reduces “ $X$ -nonnegativity” to the case  $X = \mathbb{R}^n$ .

## The standard recipe

- 1 Adopt a representation  $X = \{x : g(x) \geq 0\}$ .
- 2 Find an identity

$$f = \mathcal{L} + \sum_i \lambda_i g_i$$

where  $\mathcal{L}$  and  $\lambda_i$  are known to be nonnegative on  $\mathbb{R}^n$ .

E.g., the *positivstellensatz* of Putinar, Stengle, or Schmüdgen.

# Use the sageopt python package.



- Python 3.5 or higher (recommend  $\geq 3.6$ ).
- `"pip install sageopt"`
- Signomial and polynomial optimization.
- Require open-source convex solver, ECOS.
- Recommend commercial solver, MOSEK.