

# Signomial and Polynomial Optimization via Relative Entropy and Partial Dualization

Riley Murray

Presented at MPI MiS  
Leipzig, Germany

2 October 2019

Joint work with Venkat Chandrasekaran and Adam Wierman.

# Definitions from convex analysis

$X \subset \mathbb{R}^n$  is a **convex set** if it contains all of its line segments.

$f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a **convex function** if

$$f(tx + (1 - t)y) \leq tf(x) + (1 - t)f(y).$$

for all  $x, y \in \text{dom } f$  and all  $t \in [0, 1]$ .

# Definitions from convex analysis

$X \subset \mathbb{R}^n$  is a **convex set** if it contains all of its line segments.

$f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a **convex function** if

$$f(tx + (1 - t)y) \leq tf(x) + (1 - t)f(y).$$

for all  $x, y \in \text{dom } f$  and all  $t \in [0, 1]$ .

A convex set  $X$  induces a **support function**

$$\sigma_X(\lambda) = \sup\{\lambda^\top x : x \text{ in } X\}.$$

# Definitions from convex analysis

$X \subset \mathbb{R}^n$  is a **convex set** if it contains all of its line segments.

$f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a **convex function** if

$$f(tx + (1 - t)y) \leq tf(x) + (1 - t)f(y).$$

for all  $x, y \in \text{dom } f$  and all  $t \in [0, 1]$ .

A convex set  $X$  induces a **support function**

$$\sigma_X(\lambda) = \sup\{\lambda^T x : x \text{ in } X\}.$$

The **relative entropy function** continuously extends

$$D(u, v) = \sum_{i=1}^m u_i \log(u_i/v_i) \quad \text{to} \quad \mathbb{R}_+^m \times \mathbb{R}_+^m.$$

# Convex duality

Start with a **primal** problem

$$\text{Val}(\mathbf{c}) = \inf_{\mathbf{x}} \{ \mathbf{c}^\top \mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0} \}.$$

## Convex duality

Start with a **primal** problem

$$\text{Val}(\mathbf{c}) = \inf_{\mathbf{x}} \{ \mathbf{c}^\top \mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0} \}.$$

Obtain a **dual** problem

$$\text{Val}(\mathbf{c}) = \sup_{\boldsymbol{\mu}} \{ \mathbf{b}^\top \boldsymbol{\mu} : \mathbf{A}^\top \boldsymbol{\mu} \leq \mathbf{c} \}.$$

## Convex duality

Start with a **primal** problem

$$\text{Val}(\mathbf{c}) = \inf_{\mathbf{x}} \{ \mathbf{c}^\top \mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0} \}.$$

Obtain a **dual** problem

$$\text{Val}(\mathbf{c}) = \sup_{\boldsymbol{\mu}} \{ \mathbf{b}^\top \boldsymbol{\mu} : \mathbf{A}^\top \boldsymbol{\mu} \leq \mathbf{c} \}.$$

We will encounter constraints like

$$\text{Val}(\mathbf{c}) \geq L.$$

# Convex duality

Start with a **primal** problem

$$\text{Val}(\mathbf{c}) = \inf_{\mathbf{x}} \{ \mathbf{c}^\top \mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0} \}.$$

Obtain a **dual** problem

$$\text{Val}(\mathbf{c}) = \sup_{\boldsymbol{\mu}} \{ \mathbf{b}^\top \boldsymbol{\mu} : \mathbf{A}^\top \boldsymbol{\mu} \leq \mathbf{c} \}.$$

We will encounter constraints like

$$\text{Val}(\mathbf{c}) \geq L.$$

Write such a constraint as: *there exists* a  $\boldsymbol{\mu}$  where

$$\mathbf{A}^\top \boldsymbol{\mu} \leq \mathbf{c} \quad \text{and} \quad \mathbf{b}^\top \boldsymbol{\mu} \geq L.$$

# Nonnegativity and Optimization

Start with a function  $f$  and a set  $X \subset \mathbb{R}^n$

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : f(\mathbf{x}) \geq \gamma \text{ for all } \mathbf{x} \text{ in } X\}.$$

# Nonnegativity and Optimization

Start with a function  $f$  and a set  $X \subset \mathbb{R}^n$

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : f(\mathbf{x}) \geq \gamma \text{ for all } \mathbf{x} \text{ in } X\}.$$

**Plan of attack** for producing lower bounds:

- 1 Express  $f - \gamma = \sum_{i=1}^m c_i(\gamma) \cdot \phi_i$  for some basis functions  $\phi_i$ .

# Nonnegativity and Optimization

Start with a function  $f$  and a set  $X \subset \mathbb{R}^n$

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : f(\mathbf{x}) \geq \gamma \text{ for all } \mathbf{x} \text{ in } X\}.$$

**Plan of attack** for producing lower bounds:

- 1 Express  $f - \gamma = \sum_{i=1}^m \textcolor{blue}{c}_i(\gamma) \cdot \phi_i$  for some basis functions  $\phi_i$ .
- 2 Develop an **inner approximation**

$$\tilde{C}(\phi, X) \subset \left\{ \mathbf{c} : \sum_{i=1}^m c_i \phi_i(\mathbf{x}) \geq 0 \text{ for all } \mathbf{x} \text{ in } X \right\}.$$

# Nonnegativity and Optimization

Start with a function  $f$  and a set  $X \subset \mathbb{R}^n$

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : f(\mathbf{x}) \geq \gamma \text{ for all } \mathbf{x} \text{ in } X\}.$$

**Plan of attack** for producing lower bounds:

- 1 Express  $f - \gamma = \sum_{i=1}^m \mathbf{c}_i(\gamma) \cdot \phi_i$  for some basis functions  $\phi_i$ .
- 2 Develop an **inner approximation**

$$\tilde{C}(\phi, X) \subset \left\{ \mathbf{c} : \sum_{i=1}^m c_i \phi_i(\mathbf{x}) \geq 0 \text{ for all } \mathbf{x} \text{ in } X \right\}.$$

- 3 Find largest  $\gamma$  so  $\mathbf{c}(\gamma)$  belongs to  $\tilde{C}(\phi, X)$ .

# Our functions of interest



*polynomials*

*signomials*

# Our functions of interest

*polynomials*

*signomials*

Parameters  $\alpha_i$  in  $\mathbb{N}^n$ ,  $c_i$  in  $\mathbb{R}$ .

Using  $x^{\alpha_i} = \prod_{j=1}^n x_j^{\alpha_{ij}}$ ,

$$x \mapsto \sum_{i=1}^m c_i x^{\alpha_i}.$$

## Our functions of interest

*polynomials**signomials*Parameters  $\alpha_i$  in  $\mathbb{N}^n$ ,  $c_i$  in  $\mathbb{R}$ .Parameters  $\alpha_i$  in  $\mathbb{R}^n$ ,  $c_i$  in  $\mathbb{R}$ .Using  $x^{\alpha_i} = \prod_{j=1}^n x_j^{\alpha_{ij}}$ ,

$$x \mapsto \sum_{i=1}^m c_i x^{\alpha_i}.$$

# Our functions of interest

*polynomials*

Parameters  $\alpha_i$  in  $\mathbb{N}^n$ ,  $c_i$  in  $\mathbb{R}$ .

Using  $x^{\alpha_i} = \prod_{j=1}^n x_j^{\alpha_{ij}}$ ,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \mathbf{x}^{\alpha_i}.$$

*signomials*

Parameters  $\alpha_i$  in  $\mathbb{R}^n$ ,  $c_i$  in  $\mathbb{R}$ .

In “exponential form”,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}).$$

# Our functions of interest

## *polynomials*

Parameters  $\alpha_i$  in  $\mathbb{N}^n$ ,  $c_i$  in  $\mathbb{R}$ .

Using  $x^{\alpha_i} = \prod_{j=1}^n x_j^{\alpha_{ij}}$ ,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \mathbf{x}^{\alpha_i}.$$

Countable basis.

## *signomials*

Parameters  $\alpha_i$  in  $\mathbb{R}^n$ ,  $c_i$  in  $\mathbb{R}$ .

In “exponential form”,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}).$$

Uncountable basis.

# Our functions of interest

## *polynomials*

Parameters  $\alpha_i$  in  $\mathbb{N}^n$ ,  $c_i$  in  $\mathbb{R}$ .

Using  $x^{\alpha_i} = \prod_{j=1}^n x_j^{\alpha_{ij}}$ ,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \mathbf{x}^{\alpha_i}.$$

**Countable** basis.

Complexity measured by **degree**  
 $\max_i \alpha_{i1} + \dots + \alpha_{in}$ .

## *signomials*

Parameters  $\alpha_i$  in  $\mathbb{R}^n$ ,  $c_i$  in  $\mathbb{R}$ .

In “exponential form”,

$$\mathbf{x} \mapsto \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}).$$

**Uncountable** basis.

Complexity measured by  
**number of terms**  $m$ .

# How we think of signomials

Signomials are often written  $\mathbf{y} \mapsto \sum_{i=1}^m c_i \mathbf{y}^{\alpha_i}$ , with  $\mathbf{y} \in \mathbb{R}_{++}^n$ .

The *exponential form* has a powerful connection to convexity.

# How we think of signomials

Signomials are often written  $\mathbf{y} \mapsto \sum_{i=1}^m c_i \mathbf{y}^{\alpha_i}$ , with  $\mathbf{y} \in \mathbb{R}_{++}^n$ .

The *exponential form* has a powerful connection to convexity.

One use of your existing intuition:

- Pick an “interesting” polynomial  $p$ .
- Define  $f(z) = p(\exp(z) - \exp(-z))$ .
- $f$  will behave similarly to  $p$ , near  $0$ .

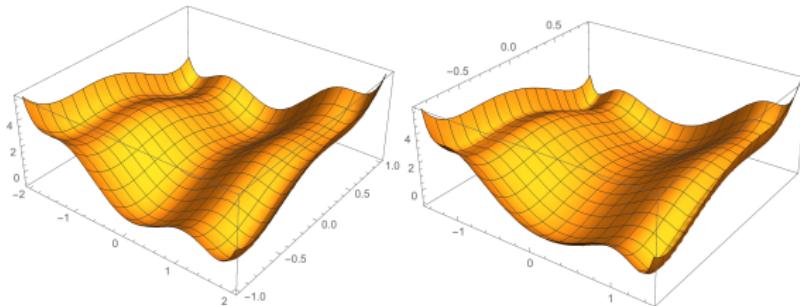
# How we think of signomials

Signomials are often written  $\mathbf{y} \mapsto \sum_{i=1}^m c_i \mathbf{y}^{\alpha_i}$ , with  $\mathbf{y} \in \mathbb{R}_{++}^n$ .

The *exponential form* has a powerful connection to convexity.

One use of your existing intuition:

- Pick an “interesting” polynomial  $p$ .
- Define  $f(z) = p(\exp(z) - \exp(-z))$ .
- $f$  will behave similarly to  $p$ , near 0.



# Outline for the talk



This talk is about **“SAGE certificates.”**

# Outline for the talk

This talk is about “**SAGE certificates**.”

## 1 SAGE signomials

Definition → Representation → Example.

# Outline for the talk

This talk is about “**SAGE certificates**.”

## 1 SAGE signomials

Definition → Representation → Example.

## 2 Signomial optimization

- Simple “SAGE relaxations.”
- Partial dualization.
- Two examples.

# Outline for the talk

This talk is about “**SAGE certificates**.”

## 1 SAGE signomials

Definition → Representation → Example.

## 2 Signomial optimization

- Simple “SAGE relaxations.”
- Partial dualization.
- Two examples.

## 3 SAGE polynomials

Definition → Example → Representation.

The signomial  $X$ -nonnegativity cones

The  $X$ -nonnegativity cone for signomials over exponents  $\alpha$ :

$$C_{\text{NNS}}(\alpha, X) \doteq \left\{ \mathbf{c} : \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}) \geq 0 \text{ for all } \mathbf{x} \text{ in } X \right\}.$$

# The signomial $X$ -nonnegativity cones

The  $X$ -nonnegativity cone for signomials over exponents  $\alpha$ :

$$C_{\text{NNS}}(\alpha, X) \doteq \left\{ \mathbf{c} : \sum_{i=1}^m c_i \exp(\alpha_i \cdot \mathbf{x}) \geq 0 \text{ for all } \mathbf{x} \text{ in } X \right\}.$$

If

$$f(\mathbf{x}) = c_1 \exp(\mathbf{0} \cdot \mathbf{x}) + c_2 \exp(\alpha_2 \cdot \mathbf{x}) + \cdots + c_m \exp(\alpha_m \cdot \mathbf{x}),$$

then

$$\inf_{\mathbf{x} \in X} f(\mathbf{x}) = \sup \{ \gamma : \mathbf{c} - (\gamma, 0, \dots, 0) \in C_{\text{NNS}}(\alpha, X) \}.$$

$$X\text{-SAGE} \Rightarrow X\text{-nonnegativity}$$


*Definition.* A signomial which is nonnegative over  $X$  and which has at most one negative coefficient is an “ $X$ -AGE function.”

$$X\text{-SAGE} \Rightarrow X\text{-nonnegativity}$$

*Definition.* A signomial which is nonnegative over  $X$  and which has at most one negative coefficient is an “ $X$ -AGE function.”

We take sums of  $X$ -AGE cones to obtain the  **$X$ -SAGE cone**

$$C_{\text{SAGE}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \underbrace{\{ \mathbf{c} : \mathbf{c}_{\setminus k} \geq \mathbf{0} \text{ and } \mathbf{c} \text{ in } C_{\text{NNS}}(\boldsymbol{\alpha}, X) \}}_{k^{\text{th}} \text{ } X\text{-AGE cone}}.$$

$$X\text{-SAGE} \Rightarrow X\text{-nonnegativity}$$

*Definition.* A signomial which is nonnegative over  $X$  and which has at most one negative coefficient is an “ $X$ -AGE function.”

We take sums of  $X$ -AGE cones to obtain the  **$X$ -SAGE cone**

$$C_{\text{SAGE}}(\alpha, X) = \sum_{k=1}^m \underbrace{\{c : c_{\setminus k} \geq \mathbf{0} \text{ and } c \text{ in } C_{\text{NNS}}(\alpha, X)\}}_{k^{\text{th}} \text{ } X\text{-AGE cone}}.$$

*Crucial question:* How to represent the AGE cones?

## The convex duality behind AGE cones

Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

## The convex duality behind AGE cones

Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

Divide out the problematic exponential, and rearrange terms:

$$\sum_{i=1}^m c_i \exp(\alpha_i \cdot x) \geq 0 \Leftrightarrow \begin{aligned} \sum_{i=1}^m c_i \exp([\alpha_i - \alpha_k] \cdot x) &\geq 0 \\ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) &\geq -c_k. \end{aligned}$$

# The convex duality behind AGE cones

Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

Divide out the problematic exponential, and rearrange terms:

$$\sum_{i=1}^m c_i \exp(\alpha_i \cdot x) \geq 0 \Leftrightarrow \begin{aligned} \sum_{i=1}^m c_i \exp([\alpha_i - \alpha_k] \cdot x) &\geq 0 \\ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) &\geq -c_k. \end{aligned}$$

The nonnegativity condition

$$\inf \left\{ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) : x \text{ in } X \right\} \geq -c_k$$

holds **if and only if**

# The convex duality behind AGE cones

Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

Divide out the problematic exponential, and rearrange terms:

$$\begin{aligned} \sum_{i=1}^m c_i \exp(\alpha_i \cdot x) \geq 0 \Leftrightarrow \sum_{i=1}^m c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq 0 \\ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq -c_k. \end{aligned}$$

The nonnegativity condition

$$\inf \left\{ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) : x \text{ in } X \right\} \geq -c_k$$

holds **if and only if** there exists  $\nu$  in  $\mathbb{R}^{m-1}$ ,  $\lambda$  in  $\mathbb{R}^n$  satisfying

# The convex duality behind AGE cones

Fix  $\alpha$  in  $\mathbb{R}^{m \times n}$ , and  $c$  in  $\mathbb{R}^m$  satisfying  $c_{\setminus k} \geq 0$ . Convex  $X \subset \mathbb{R}^n$ .

Divide out the problematic exponential, and rearrange terms:

$$\begin{aligned} \sum_{i=1}^m c_i \exp(\alpha_i \cdot x) \geq 0 \Leftrightarrow \sum_{i=1}^m c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq 0 \\ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) \geq -c_k. \end{aligned}$$

The nonnegativity condition

$$\inf \left\{ \sum_{i \neq k} c_i \exp([\alpha_i - \alpha_k] \cdot x) : x \text{ in } X \right\} \geq -c_k$$

holds **if and only if** there exists  $\nu$  in  $\mathbb{R}^{m-1}$ ,  $\lambda$  in  $\mathbb{R}^n$  satisfying

$$\begin{aligned} [\mathbf{1}\alpha_k - \alpha_{\setminus k}]^\top \nu = \lambda \quad \text{and} \\ \sigma_X(\lambda) + D(\nu, c_{\setminus k}) - \nu^\top \mathbf{1} \leq c_k. \end{aligned}$$

Tractability of  $X$ -SAGE cones

There are two constraints in an AGE cone:

- $[\mathbf{1}\alpha_k - \alpha_{\setminus k}]^\top \boldsymbol{\nu} = \boldsymbol{\lambda}$ , and
- $\sigma_X(\boldsymbol{\lambda}) + D(\boldsymbol{\nu}, \mathbf{c}_{\setminus k}) - \boldsymbol{\nu}^\top \mathbf{1} \leq c_k$ .

Support function  $\sigma_X(\boldsymbol{\lambda}) = \sup\{\boldsymbol{\lambda}^\top \mathbf{x} : \mathbf{x} \text{ in } X\}$  always convex.

Tractability of  $X$ -SAGE cones

There are two constraints in an AGE cone:

- $[\mathbf{1}\boldsymbol{\alpha}_k - \boldsymbol{\alpha}_{\setminus k}]^\top \boldsymbol{\nu} = \boldsymbol{\lambda}$ , and
- $\sigma_X(\boldsymbol{\lambda}) + D(\boldsymbol{\nu}, \mathbf{c}_{\setminus k}) - \boldsymbol{\nu}^\top \mathbf{1} \leq c_k$ .

Support function  $\sigma_X(\boldsymbol{\lambda}) = \sup\{\boldsymbol{\lambda}^\top \mathbf{x} : \mathbf{x} \text{ in } X\}$  always convex.

When does it have a closed form? Examples include ...

$$X = \mathbb{R}^n \quad \Rightarrow \quad \sigma_X(\boldsymbol{\lambda}) = \begin{cases} 0 & \text{if } \boldsymbol{\lambda} = \mathbf{0} \\ +\infty & \text{if } \boldsymbol{\lambda} \neq \mathbf{0} \end{cases}$$

Tractability of  $X$ -SAGE cones

There are two constraints in an AGE cone:

- $[\mathbf{1}\alpha_k - \alpha_{\setminus k}]^\top \boldsymbol{\nu} = \boldsymbol{\lambda}$ , and
- $\sigma_X(\boldsymbol{\lambda}) + D(\boldsymbol{\nu}, \mathbf{c}_{\setminus k}) - \boldsymbol{\nu}^\top \mathbf{1} \leq c_k$ .

Support function  $\sigma_X(\boldsymbol{\lambda}) = \sup\{\boldsymbol{\lambda}^\top \mathbf{x} : \mathbf{x} \text{ in } X\}$  always convex.

When does it have a closed form? Examples include ...

$$X = \{\mathbf{x} : \|\mathbf{x} - \mathbf{a}\| \leq r\} \quad \Rightarrow \quad \sigma_X(\boldsymbol{\lambda}) = \boldsymbol{\lambda}^\top \mathbf{a} + r\|\boldsymbol{\lambda}\|_*$$

# Tractability of $X$ -SAGE cones

There are two constraints in an AGE cone:

- $[\mathbf{1}\alpha_k - \alpha_{\setminus k}]^\top \boldsymbol{\nu} = \boldsymbol{\lambda}$ , and
- $\sigma_X(\boldsymbol{\lambda}) + D(\boldsymbol{\nu}, \mathbf{c}_{\setminus k}) - \boldsymbol{\nu}^\top \mathbf{1} \leq c_k$ .

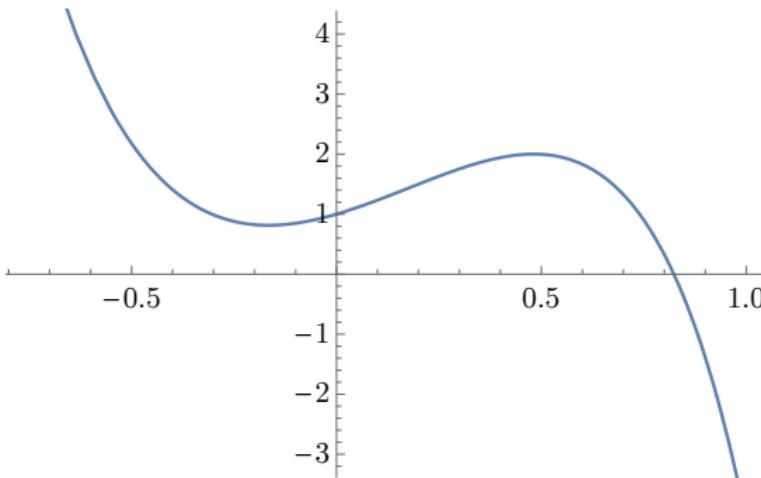
Support function  $\sigma_X(\boldsymbol{\lambda}) = \sup\{\boldsymbol{\lambda}^\top \mathbf{x} : \mathbf{x} \text{ in } X\}$  always convex.

When is it *tractable*?

**Whenever  $X$  is tractable.**

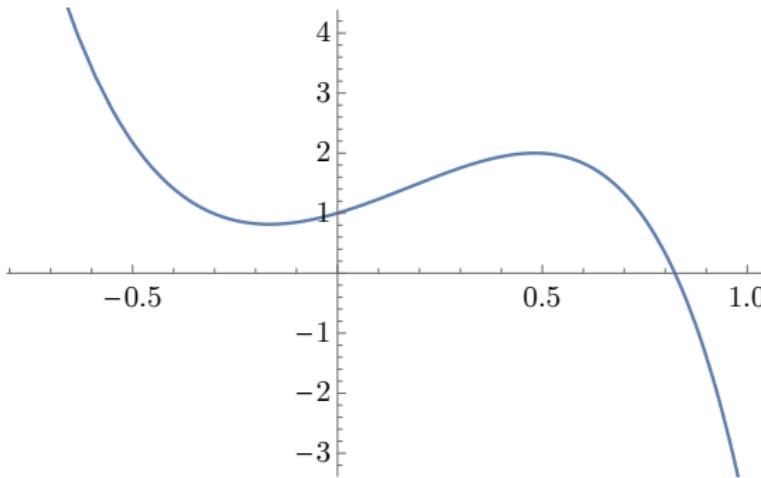
# A univariate example

$$f(x) = e^{-3x} + e^{-2x} + 4e^x + e^{2x} - 4e^{-x} - 1 - e^{3x} \text{ over } x \leq 0$$



# A univariate example

$$f(x) = e^{-3x} + e^{-2x} + 4e^x + e^{2x} - 4e^{-x} - 1 - e^{3x} \text{ over } x \leq 0$$



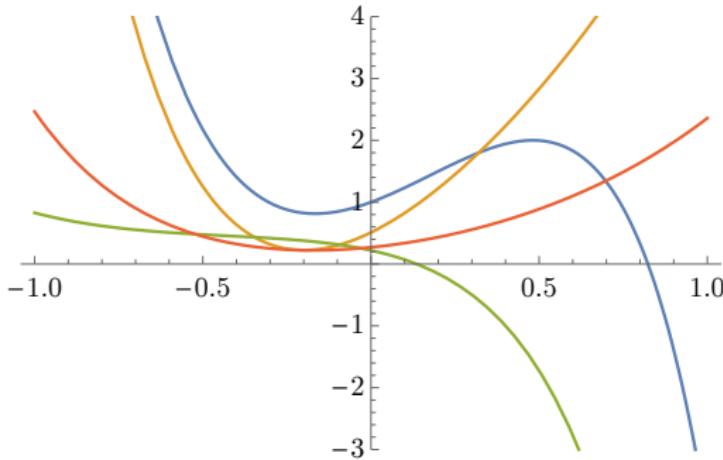
$$f_1(x) = 0.88 \cdot e^{-3x} + 0.82 \cdot e^{-2x} + 2.69 \cdot e^x + 0.12 \cdot e^{2x} - 4 \cdot e^{-x}$$

$$f_2(x) = 0.10 \cdot e^{-3x} + 0.15 \cdot e^{-2x} + 0.90 \cdot e^x + 0.12 \cdot e^{2x} - 1$$

$$f_3(x) = 0.02 \cdot e^{-3x} + 0.03 \cdot e^{-2x} + 0.41 \cdot e^x + 0.76 \cdot e^{2x} - e^{3x}$$

# A univariate example

$$f(x) = e^{-3x} + e^{-2x} + 4e^x + e^{2x} - 4e^{-x} - 1 - e^{3x} \text{ over } x \leq 0$$



$$f_1(x) = 0.88 \cdot e^{-3x} + 0.82 \cdot e^{-2x} + 2.69 \cdot e^x + 0.12 \cdot e^{2x} - 4 \cdot e^{-x}$$

$$f_2(x) = 0.10 \cdot e^{-3x} + 0.15 \cdot e^{-2x} + 0.90 \cdot e^x + 0.12 \cdot e^{2x} - 1$$

$$f_3(x) = 0.02 \cdot e^{-3x} + 0.03 \cdot e^{-2x} + 0.41 \cdot e^x + 0.76 \cdot e^{2x} - e^{3x}$$

# Optimization.

## Simple SAGE relaxations

Consider  $f(\mathbf{x}) = \sum_{i=1}^m c_i \exp(\boldsymbol{\alpha}_i \cdot \mathbf{x})$  with  $\boldsymbol{\alpha}_1 = \mathbf{0}$ .

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{NNS}}(\boldsymbol{\alpha}, X)\}$$

$$\geq \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)\}$$

$$= \inf \left\{ \mathbf{c}^\top \mathbf{v} : \begin{array}{l} \mathbf{v} \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)^\dagger \\ \text{satisfies } \mathbf{v} \cdot \mathbf{e}_1 = 1 \end{array} \right\}$$

What about [solution recovery](#)?

## Simple SAGE relaxations

Consider  $f(\mathbf{x}) = \sum_{i=1}^m c_i \exp(\boldsymbol{\alpha}_i \cdot \mathbf{x})$  with  $\boldsymbol{\alpha}_1 = \mathbf{0}$ .

$$\inf\{f(\mathbf{x}) : \mathbf{x} \text{ in } X\} = \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{NNS}}(\boldsymbol{\alpha}, X)\}$$

$$\geq \sup\{\gamma : \mathbf{c} - \gamma \mathbf{e}_1 \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)\}$$

$$= \inf \left\{ \mathbf{c}^\top \mathbf{v} : \begin{array}{l} \mathbf{v} \text{ in } C_{\text{SAGE}}(\boldsymbol{\alpha}, X)^\dagger \\ \text{satisfies } \mathbf{v} \cdot \mathbf{e}_1 = 1 \end{array} \right\}$$

What about **solution recovery**? When  $X$  is convex, we have

$$\begin{aligned} C_{\text{SAGE}}(\boldsymbol{\alpha}, X)^\dagger = \text{cl}\{ \mathbf{v} : \text{some } \mathbf{z}_1, \dots, \mathbf{z}_m \text{ in } \mathbb{R}^n \text{ satisfy} \\ v_k \log(\mathbf{v}/v_k) \geq [\boldsymbol{\alpha} - \mathbf{1}\boldsymbol{\alpha}_k] \mathbf{z}_k \\ \text{and } \mathbf{z}_k/v_k \in X \text{ for all } k \text{ in } [m] \}. \end{aligned}$$

An example in  $\mathbb{R}^3$ 

Minimize

$$f(\mathbf{x}) = 0.5 \exp(x_1 - x_2) - \exp x_1 - 5 \exp(-x_2)$$

over

$$\begin{aligned} X = \{ \mathbf{x} : \log 70 \leq x_1 \leq \log 150, \\ \log 1.0 \leq x_2 \leq \log 30, \\ \log 0.5 \leq x_3 \leq \log 21 \\ \exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100 \}. \end{aligned}$$

An example in  $\mathbb{R}^3$ 

Minimize

$$f(\mathbf{x}) = 0.5 \exp(x_1 - x_2) - \exp x_1 - 5 \exp(-x_2)$$

over

$$\begin{aligned} X = \{ \mathbf{x} : \log 70 \leq x_1 \leq \log 150, \\ \log 1.0 \leq x_2 \leq \log 30, \\ \log 0.5 \leq x_3 \leq \log 21 \\ \exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100 \}. \end{aligned}$$

Compute  $f_X^{\text{SAGE}} = -147.85713 \leq f_X^{\star}$ ,

An example in  $\mathbb{R}^3$ 

Minimize

$$f(\mathbf{x}) = 0.5 \exp(x_1 - x_2) - \exp x_1 - 5 \exp(-x_2)$$

over

$$\begin{aligned} X = \{ \mathbf{x} : \log 70 \leq x_1 \leq \log 150, \\ \log 1.0 \leq x_2 \leq \log 30, \\ \log 0.5 \leq x_3 \leq \log 21 \\ \exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100 \}. \end{aligned}$$

Compute  $f_X^{\text{SAGE}} = -147.85713 \leq f_X^*$ , recover feasible

$$\mathbf{x}^* = (5.01063529, 3.40119660, -0.48450710)$$

satisfying  $f(\mathbf{x}^*) = -147.66666$ .

An example in  $\mathbb{R}^3$ 

Minimize

$$f(\mathbf{x}) = 0.5 \exp(x_1 - x_2) - \exp x_1 - 5 \exp(-x_2)$$

over

$$\begin{aligned} X = \{ \mathbf{x} : \log 70 \leq x_1 \leq \log 150, \\ \log 1.0 \leq x_2 \leq \log 30, \\ \log 0.5 \leq x_3 \leq \log 21 \\ \exp(x_2 - x_3) + \exp x_2 + 0.05 \exp(x_1 + x_3) \leq 100 \}. \end{aligned}$$

Compute  $f_X^{\text{SAGE}} = -147.85713 \leq f_X^*$ , recover feasible

$$\mathbf{x}^* = (5.01063529, 3.40119660, -0.48450710)$$

satisfying  $f(\mathbf{x}^*) = -147.66666$ . *This is actually optimal!*

# Nonconvex constraints



Q: What should we do when some constraints are nonconvex?

# Nonconvex constraints



Q: What should we do when some constraints are nonconvex?

A: Combine  $X$ -SAGE certificates with Lagrangian relaxations.

# Nonconvex constraints

Q: What should we do when some constraints are nonconvex?

A: Combine  $X$ -SAGE certificates with Lagrangian relaxations.

Concretely, suppose we want to minimize  $f$  over

$$\Omega \doteq X \cap \{\mathbf{x} : g(\mathbf{x}) \leq \mathbf{0}\}$$

where  $X$  is convex, but  $g_1, \dots, g_k$  are nonconvex signomials.

# Nonconvex constraints

Q: What should we do when some constraints are nonconvex?

A: Combine  $X$ -SAGE certificates with Lagrangian relaxations.

Concretely, suppose we want to minimize  $f$  over

$$\Omega \doteq X \cap \{\mathbf{x} : g(\mathbf{x}) \leq \mathbf{0}\}$$

where  $X$  is convex, but  $g_1, \dots, g_k$  are nonconvex signomials.

Then, if  $\lambda_1, \dots, \lambda_k$  are nonnegative dual variables, we have

$$\inf_{\mathbf{x} \in \Omega} f(\mathbf{x}) \geq \sup \left\{ \gamma : f + \sum_{i=1}^k \lambda_i g_i - \gamma \text{ is } X\text{-SAGE} \right\}.$$

# The SimPleAC aircraft design problem

From Warren Hoburg's PhD thesis.

Problem statistics:

- 140 variables.
- 89 inequality constraints (1 nonconvex).
- 67 equality constraints (15 nonconvex).

Performance of the most basic SAGE relaxation:

- bound “ $\text{cost} \geq 2957$ ” (roughly match a known solution).
- MOSEK solves in two seconds, on a six year old laptop.
- solution recovery fails (numerical issues).

# SAGE polynomials

# $X$ -nonnegative and $X$ -AGE polynomials

Recall our standard notation  $\mathbf{x}^{\alpha_i} \doteq \prod_{j=1}^n x_j^{\alpha_{ij}}$ .

The matrix  $\alpha$  and a set  $X \subset \mathbb{R}^n$  induce a nonnegativity cone

$$C_{\text{NNP}}(\alpha, X) = \{ \mathbf{c} : c_1 \mathbf{x}^{\alpha_1} + \cdots + c_m \mathbf{x}^{\alpha_m} \geq 0 \text{ for all } \mathbf{x} \text{ in } X \}.$$

# $X$ -nonnegative and $X$ -AGE polynomials

Recall our standard notation  $\mathbf{x}^{\alpha_i} \doteq \prod_{j=1}^n x_j^{\alpha_{ij}}$ .

The matrix  $\alpha$  and a set  $X \subset \mathbb{R}^n$  induce a nonnegativity cone

$$C_{\text{NNP}}(\alpha, X) = \{ \mathbf{c} : c_1 \mathbf{x}^{\alpha_1} + \cdots + c_m \mathbf{x}^{\alpha_m} \geq 0 \text{ for all } \mathbf{x} \text{ in } X \}.$$

Def.  $f(\mathbf{x}) = c_1 \mathbf{x}^{\alpha_1} + \cdots + c_m \mathbf{x}^{\alpha_m}$  is an  **$X$ -AGE polynomial** if

- 1  $\mathbf{c}$  belongs to  $C_{\text{NNP}}(\alpha, X)$ , and
- 2 at most one “ $k$ ” has  $c_k \mathbf{x}^{\alpha_k} < 0$  for some  $\mathbf{x} \in X$ .

## A symbolic example

Consider  $f(\mathbf{x}) = n - \sum_{i=1}^n \prod_{j \in [n] \setminus \{i\}} x_j$  over  $X = \{-1, 1\}^n$ .

Let's prove that  $f$  is  $X$ -nonnegative.

## A symbolic example

Consider  $f(\mathbf{x}) = n - \sum_{i=1}^n \prod_{j \in [n] \setminus \{i\}} x_j$  over  $X = \{-1, 1\}^n$ .

Let's prove that  $f$  is  $X$ -nonnegative.

- 1 Note that  $\mathbf{x} \in \{-1, 1\}^n$  implies  $\prod_{j \neq i} x_j \leq 1$ .

## A symbolic example

Consider  $f(\mathbf{x}) = n - \sum_{i=1}^n \prod_{j \in [n] \setminus \{i\}} x_j$  over  $X = \{-1, 1\}^n$ .

Let's prove that  $f$  is  $X$ -nonnegative.

- 1 Note that  $\mathbf{x} \in \{-1, 1\}^n$  implies  $\prod_{j \neq i} x_j \leq 1$ .
- 2 Therefore  $f_i(\mathbf{x}) = 1 - \prod_{j \neq i} x_j$  are  $X$ -AGE.

## A symbolic example

Consider  $f(\mathbf{x}) = n - \sum_{i=1}^n \prod_{j \in [n] \setminus \{i\}} x_j$  over  $X = \{-1, 1\}^n$ .

Let's prove that  $f$  is  $X$ -nonnegative.

- 1 Note that  $\mathbf{x} \in \{-1, 1\}^n$  implies  $\prod_{j \neq i} x_j \leq 1$ .
- 2 Therefore  $f_i(\mathbf{x}) = 1 - \prod_{j \neq i} x_j$  are  $X$ -AGE.
- 3 Since  $f = \sum_{i=1}^n f_i$ , conclude  $f$  is  $X$ -SAGE.

# The representation problem

The cone of coefficients for  $X$ -**SAGE polynomials** is given by

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \mathbf{c} \mid \begin{array}{l} \mathbf{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \mathbf{x} \in X \Rightarrow c_i \mathbf{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}.$$

# The representation problem

The cone of coefficients for  $X$ -**SAGE polynomials** is given by

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \boldsymbol{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \boldsymbol{x} \in X \Rightarrow c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}.$$

- This definition cares about the behavior of monomials.

# The representation problem

The cone of coefficients for  $X$ -**SAGE polynomials** is given by

$$C_{\text{SAGE}}^{\text{POLY}}(\alpha, X) = \sum_{k=1}^m \left\{ \mathbf{c} \mid \begin{array}{l} \mathbf{c} \text{ in } C_{\text{NNP}}(\alpha, X), \text{ and for} \\ i \neq k, \mathbf{x} \in X \Rightarrow c_i \mathbf{x}^{\alpha_i} \geq 0 \end{array} \right\}.$$

- This definition cares about the behavior of monomials.
- SAGE signomials put monomials front and center.

# The representation problem

The cone of coefficients for  $X$ -**SAGE polynomials** is given by

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \mathbf{c} \mid \begin{array}{l} \mathbf{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \mathbf{x} \in X \Rightarrow c_i \mathbf{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}.$$

- This definition cares about the behavior of monomials.
- SAGE signomials put monomials front and center.
- Monomials " $\mathbf{x}^{\boldsymbol{\alpha}_i}$ " are quite different from " $\exp(\boldsymbol{\alpha}_i \cdot \mathbf{x})$ ".

## Representation: the single-orthant case

Consider  $X \subset \mathbb{R}_+^n$ , and define  $X_{++} \doteq X \cap \mathbb{R}_{++}^n$ .

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \boldsymbol{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \boldsymbol{x} \in X \Rightarrow c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}$$

## Representation: the single-orthant case

Consider  $X \subset \mathbb{R}_+^n$ , and define  $X_{++} \doteq X \cap \mathbb{R}_{++}^n$ .

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \boldsymbol{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \boldsymbol{x} \in X \Rightarrow c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}$$

**Fact 1.** If  $X = \text{cl } X_{++}$ , then  $C_{\text{NNP}}(\boldsymbol{\alpha}, X) = C_{\text{NNS}}(\boldsymbol{\alpha}, \log X_{++})$ .

## Representation: the single-orthant case

Consider  $X \subset \mathbb{R}_+^n$ , and define  $X_{++} \doteq X \cap \mathbb{R}_{++}^n$ .

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \text{c in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \boldsymbol{x} \in X \Rightarrow c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}$$

**Fact 1.** If  $X = \text{cl } X_{++}$ , then  $C_{\text{NNP}}(\boldsymbol{\alpha}, X) = C_{\text{NNS}}(\boldsymbol{\alpha}, \log X_{++})$ .

**Fact 2.** Since  $X \subset \mathbb{R}_+^n$ , “ $c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \ \forall \boldsymbol{x} \in X$ ” reduces to  $c_i \geq 0$ .

## Representation: the single-orthant case

Consider  $X \subset \mathbb{R}_+^n$ , and define  $X_{++} \doteq X \cap \mathbb{R}_{++}^n$ .

$$C_{\text{SAGE}}^{\text{POLY}}(\boldsymbol{\alpha}, X) = \sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \boldsymbol{c} \text{ in } C_{\text{NNP}}(\boldsymbol{\alpha}, X), \text{ and for} \\ i \neq k, \boldsymbol{x} \in X \Rightarrow c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \end{array} \right\}$$

**Fact 1.** If  $X = \text{cl } X_{++}$ , then  $C_{\text{NNP}}(\boldsymbol{\alpha}, X) = C_{\text{NNS}}(\boldsymbol{\alpha}, \log X_{++})$ .

**Fact 2.** Since  $X \subset \mathbb{R}_+^n$ , “ $c_i \boldsymbol{x}^{\boldsymbol{\alpha}_i} \geq 0 \ \forall \boldsymbol{x} \in X$ ” reduces to  $c_i \geq 0$ .

$$\sum_{k=1}^m \left\{ \boldsymbol{c} \mid \begin{array}{l} \boldsymbol{c} \text{ in } C_{\text{NNS}}(\boldsymbol{\alpha}, \log X_{++}), \\ \text{and } c_i \geq 0 \text{ for all } i \neq k \end{array} \right\} = C_{\text{SAGE}}(\boldsymbol{\alpha}, \log X_{++})$$

# Representation: sign-symmetry

Consider  $X \subset \mathbb{R}^n$  which satisfies

- 1 invariance under reflection about  $\{x : x_i = 0\}$ ,
- 2 and  $X \cap \mathbb{R}_+^n = \text{cl } X_{++}$ .

For example,  $X = \{-1, 1\}^n$ .

# Representation: sign-symmetry

Consider  $X \subset \mathbb{R}^n$  which satisfies

- 1 invariance under reflection about  $\{x : x_i = 0\}$ ,
- 2 and  $X \cap \mathbb{R}_+^n = \text{cl } X_{++}$ .

For example,  $X = \{-1, 1\}^n$ .

It can subsequently be shown that

$C_{\text{SAGE}}^{\text{POLY}}(\alpha, X) = \{c : \text{some } \hat{c} \in C_{\text{SAGE}}(\alpha, \log X_{++}) \text{ satisfies}$

$$\hat{c}_i = c_i \text{ whenever } \alpha_i \text{ is in } 2\mathbb{N}^n, \text{ and}$$

$$\hat{c}_i \leq -|c_i| \text{ whenever } \alpha_i \text{ is not in } 2\mathbb{N}^n\}.$$

# Representation: sign-symmetry

Consider  $X \subset \mathbb{R}^n$  which satisfies

- 1 invariance under reflection about  $\{x : x_i = 0\}$ ,
- 2 and  $X \cap \mathbb{R}_+^n = \text{cl } X_{++}$ .

For example,  $X = \{-1, 1\}^n$ .

It can subsequently be shown that

$C_{\text{SAGE}}^{\text{POLY}}(\alpha, X) = \{c : \text{some } \hat{c} \in C_{\text{SAGE}}(\alpha, \log X_{++}) \text{ satisfies}$   
 $\hat{c}_i = c_i \text{ whenever } \alpha_i \text{ is in } 2\mathbb{N}^n, \text{ and}$   
 $\hat{c}_i \leq -|c_i| \text{ whenever } \alpha_i \text{ is not in } 2\mathbb{N}^n\}$ .

If  $\alpha_i$  isn't even, then some  $x_1, x_2 \in X$  satisfy  $x_1^{\alpha_i} < 0 < x_2^{\alpha_i}$ .

# Log-log convexity



For what  $X \subset \mathbb{R}_{++}^n$  is  $\log X$  convex?

# Log-log convexity

For what  $X \subset \mathbb{R}_{++}^n$  is  $\log X$  convex?

A function  $g$  is **log-log convex** on  $D \subset \mathbb{R}_{++}^n$  when

- $\log D$  is a convex set, and
- $\log(x) \mapsto \log g(x)$  is convex function.

Such functions are sometimes called *geometrically convex*.

# Log-log convexity

For what  $X \subset \mathbb{R}_{++}^n$  is  $\log X$  convex?

A function  $g$  is **log-log convex** on  $D \subset \mathbb{R}_{++}^n$  when

- $\log D$  is a convex set, and
- $\log(x) \mapsto \log g(x)$  is convex function.

Such functions are sometimes called *geometrically convex*.

Studied by Montel (1928) and Niculescu (2000), among others.

If  $g_1, \dots, g_k$  are log-log convex on a box  $B \subset \mathbb{R}_{++}^n$ , then

$\log\{(x, \mathbf{t}) : x \in B, \mathbf{t} \in \mathbb{R}_{++}^k, g(x) \leq \mathbf{t}\} \subset \mathbb{R}^{n+k}$  is convex.

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

$$g(x) = \left( \int_x^{\infty} e^{-t^2} dt \right)^{-1}$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

With more restricted domains:

$$x \mapsto (-x \log x)^{-1} \quad D = (0, 1)$$

$$g(x) = \left( \int_x^{\infty} e^{-t^2} dt \right)^{-1}$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

$$g(x) = \left( \int_x^{\infty} e^{-t^2} dt \right)^{-1}$$

With more restricted domains:

$$x \mapsto (-x \log x)^{-1} \quad D = (0, 1)$$

$$\mathbf{X} \mapsto (\mathbf{I} - \mathbf{X})^{-1}$$

$$D = \{\mathbf{X} \in \mathbb{R}_{++}^{n \times n} : \rho(\mathbf{X}) < 1\}$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

$$g(x) = \left( \int_x^\infty e^{-t^2} dt \right)^{-1}$$

With more restricted domains:

$$x \mapsto (-x \log x)^{-1} \quad D = (0, 1)$$

$$\mathbf{X} \mapsto (\mathbf{I} - \mathbf{X})^{-1}$$

$$D = \{\mathbf{X} \in \mathbb{R}_{++}^{n \times n} : \rho(\mathbf{X}) < 1\}$$

$$x \mapsto (\log x)^{-1} \quad D = (1, \infty)$$

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

$$g(x) = \left( \int_x^\infty e^{-t^2} dt \right)^{-1}$$

With more restricted domains:

$$x \mapsto (-x \log x)^{-1} \quad D = (0, 1)$$

$$\mathbf{X} \mapsto (\mathbf{I} - \mathbf{X})^{-1}$$

$$D = \{\mathbf{X} \in \mathbb{R}_{++}^{n \times n} : \rho(\mathbf{X}) < 1\}$$

$$x \mapsto (\log x)^{-1} \quad D = (1, \infty)$$

Some tractable constraints for  $X$ -SAGE polynomials:

## Log-log convexity: examples

With domains  $D = \mathbb{R}_{++}^n$ :

$$g(\mathbf{x}) = \max\{x_1, \dots, x_n\}$$

$$g(\mathbf{x}) = x_1^{a_1} \cdots x_n^{a_n}$$

$$g(x) = \left( \int_x^\infty e^{-t^2} dt \right)^{-1}$$

With more restricted domains:

$$x \mapsto (-x \log x)^{-1} \quad D = (0, 1)$$

$$\mathbf{X} \mapsto (\mathbf{I} - \mathbf{X})^{-1}$$

$$D = \{\mathbf{X} \in \mathbb{R}_{++}^{n \times n} : \rho(\mathbf{X}) < 1\}$$

$$x \mapsto (\log x)^{-1} \quad D = (1, \infty)$$

Some tractable constraints for  $X$ -SAGE polynomials:

$$\|\mathbf{x}\|_p \leq a \quad x_j^2 = a \quad a \leq \mathbb{P}\{\mathcal{N}(0, \sigma) \geq |x|\}$$

where  $a > 0$ .

# Thank you!

# Handling conditional nonnegativity

Typically, one reduces “ $X$ -nonnegativity” to the case  $X = \mathbb{R}^n$ .

## The standard recipe

- 1 Adopt a representation  $X = \{x : g(x) \geq \mathbf{0}\}$ .
- 2 Find an identity

$$f = \mathcal{L} + \sum_i \lambda_i g_i$$

where  $\mathcal{L}$  and  $\lambda_i$  are known to be nonnegative on  $\mathbb{R}^n$ .

E.g., the *positivstellensatz* of Putinar, Stengle, or Schmudgen.

Use the sageopt python package.



- Python 3.5 or higher (recommend  $\geq 3.6$ ).
- “`pip install sageopt`”
- Signomial and polynomial optimization.
- Require open-source convex solver, ECOS.
- Recommend commercial solver, MOSEK.